

GAME: Arabian Nights [[https://en.wikipedia.org/wiki/Arabian_Nights_\(2001_video_game\)](https://en.wikipedia.org/wiki/Arabian_Nights_(2001_video_game))]

Protection: SecuROM *new* 4.48.00.0004

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You will need:

- Windows XP VM (I used VMware [<https://www.vmware.com/products/workstation-player.html>])
- x64dbg (x32dbg) [<https://x64dbg.com/>]
- CFF Explorer [https://ntcore.com/?page_id=388]
- Original game disc (you need the ORIGINAL, otherwise this will not work)

Before you start:

SecuROM protected games may not work properly on Windows versions newer than XP.

As we already experienced with SafeDisc, once we removed this DRM, the game works perfectly even on Windows 11.

This DRM works by replacing (proxying) various Windows APIs used by the game, with a function that after running some checks, will reach the requested API with a jump instruction. This jump will be absolute, without passing from the relative *IAT* thunk, so when we rebuild the imports, we will need to loop through the *IAT* to find the correct thunk, get its address, and then replace the SecuROM call in the **.text** segment with it.

There's also an initial layer of encryption (we will need the original game disc to decrypt it) and various anti-debugging techniques that will make it harder to reach the Original Entry Point (*OEP*).

Let's begin:

Install the game selecting the *FULL* install option. Once installed load the main executable (*_start.exe*) inside the debugger.

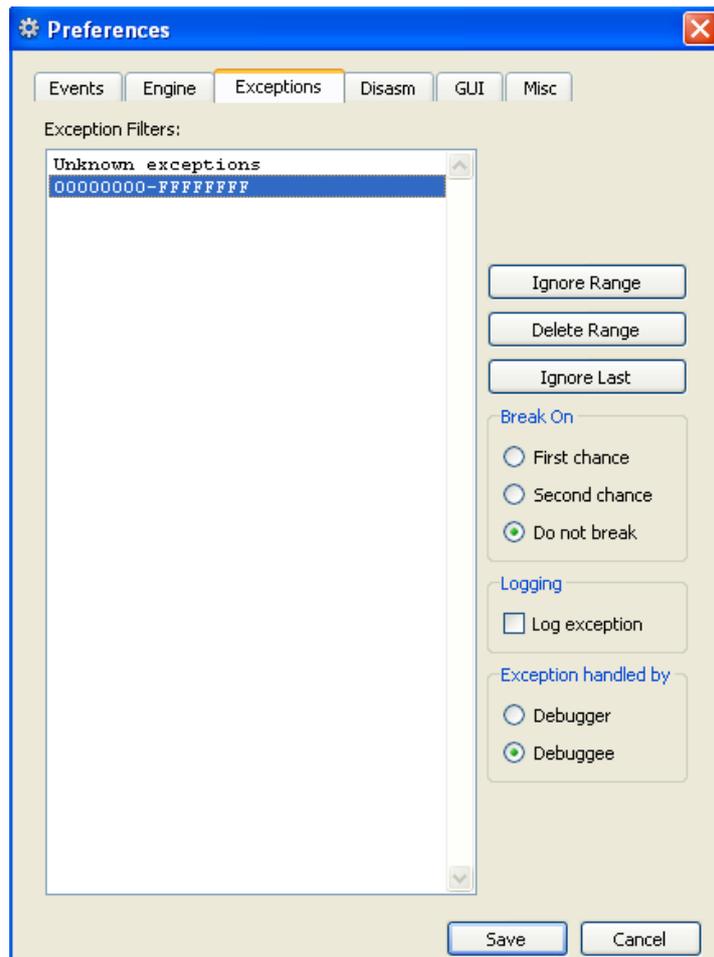
We can see that the entry point is located at **0x737CFD**.

If we go to the *Memory Map* tab, we can see that we are currently located in the **.cms_t** section:

00400000	00001000	_start.exe				
00401000	00062000	".text"	Executable code	IMG	-R---	ERWC-
00463000	00003000	".rdata"	Read-only initialized data	IMG	-R---	ERWC-
00466000	002C4000	".data"	Initialized data	IMG	-RWC-	ERWC-
0072A000	00002000	".idata"	Import tables	IMG	-RWC-	ERWC-
0072C000	00014000	".cms_t"		IMG	ER---	ERWC-
00740000	0002C000	".cms_d"		IMG	-RWC-	ERWC-
0076C000	00001000	".idata"	Import tables	IMG	-RW--	ERWC-
0076D000	00001000	".rsrc"	Resources	IMG	-R---	ERWC-
0076E000	00009000	".reloc"	Base relocations	IMG	-R---	ERWC-

We can suppose that the code of our game is in the **.text** segment and what we are going to execute is the SecuROM loader.

If we try to click on *RUN*, we will be constantly blocked with exceptions of various types: this is only one of the various techniques put in place to slow us down. This problem can be easily fixed by configuring the debugger to ignore all exceptions:



I also recommend unchecking “*Log exception*”, because the huge number of exceptions will extremely slow down the execution. Now we are ready to start.

As in most cases when trying to remove a protection of this type, the first step is to be able to reach the *OEP*.

My first attempt was to set a hardware breakpoint on the `.text` segment on execution: unfortunately thanks to the various anti-debugging techniques used, this operation will cause an endless loop of the SecuROM loader.

So, I decided to proceed in two steps:

- 1) Find out at what address the *OEP* is located
- 2) Find a way to reach said address

To get closer to the *OEP*, I’ve set a breakpoint on an API that is usually located near the Entry Point: **GetCommandLineA**.

Each time this breakpoint is triggered, we must click on “*Run to user code*” to see where the call originates from. Once we hit the breakpoint for the 3rd time, we are finally located inside the `.text` segment:

0044D033	83C4 04	add esp,4
0044D036	85C0	test eax,eax
0044D038	75 0A	jne _start.44D044
0044D03A	6A 1C	push 1C
0044D03C	E8 FF000000	call _start.44D0E40
0044D041	83C4 04	add esp,4
0044D044	C745 FC 00000000	mov dword ptr ss:[ebp-4],0
0044D048	E8 207B0000	call _start.455870
0044D050	FF15 00137400	call dword ptr ds:[741300]
0044D056	A3 18917200	mov dword ptr ds:[729118],eax
0044D058	E8 F0780000	call _start.455650
0044D060	A3 30CE4700	mov dword ptr ds:[47CE30],eax
0044D065	E8 06730000	call _start.455140
0044D06A	E8 81720000	call _start.454FF0
0044D06F	E8 5C1E0000	call _start.44FB00
0044D074	C745 00 00000000	mov dword ptr ss:[ebp-30],0

If you look closer, you will realize that the **GetCommandLineA** call was originated from **call dword ptr ds:[741300]**.

This is quite interesting, but for the moment let's focus on our current target (reaching the OEP). Since we assumed that the **GetCommandLineA** API is in the function where the OEP resides, we can scroll up a little bit till the start of the current function, and we will finally be at our destination:

0044DCB6	CC	int3
0044DCAC	CC	int3
0044DCAD	CC	int3
0044DCAE	CC	int3
0044DCAF	CC	int3
0044DCB0	55	push ebp
0044DCB1	8BEC	mov ebp,esp
0044DCB3	6A FF	push FFFFFFFF
0044DCB5	68 70454600	push _start.464570
0044DCBA	68 D4AE4400	push _start.44AED4
0044DCBF	64:A1 00000000	mov eax,dword ptr fs:[0]
0044DCC5	50	push eax
0044DCC6	64:8925 00000000	mov dword ptr fs:[0],esp
0044DCCD	83C4 A4	add esp,FFFFFFA4
0044DCD0	53	push ebx
0044DCD1	56	push esi

Perfect, now we know that **0x44DCB0** is the *OEP*. We need a way to halt the code execution right there.

PAY ATTENTION: it's mandatory to break the *OEP* when dumping from memory! Otherwise, the resulting executable will not work, as it will contain data related to the current execution.

If we try to set a breakpoint at that address, after restarting the debugger, the SecuROM loader will detect it and will cause an endless loop. No matter if the breakpoint is hardware-based: it will be still detected.

If we restart the debugger once more and we go to the address where the *OEP* resides BEFORE running the SecuROM loader, we will notice a very interesting thing:

0044DCAB	1863 31	sbb byte ptr ds:[ebx+31],ah
0044DCAE	E5 AB	in eax,AB
0044DCB0	6A E6	push FFFFFFFE6
0044DCB2	46	inc esi
0044DCB3	72 D7	jb _start.44DC8C
0044DCB5	2B30	sub esi,dword ptr ds:[eax]
0044DCB7	E5 28	in eax,28
0044DCB9	B8 6B49CC35	mov eax,35CC496B
0044DCBE	FE	ret
0044DCBF	BF F040CFBA	mov edi,BACF40F0
0044DCC4	3A65 4E	cmp ah,byte ptr ss:[ebp+4E]
0044DCC7	C8 44D3 20	enter 0344,20
0044DCCB	92	xchg edx,eax

This code is completely different from what we expected to see! It is reasonable to think that this is encrypted memory and that it will be overwritten by the SecuROM loader during the boot phase.

On Windows, processes can modify memory thanks to the **WriteProcessMemory** API, which has the following signature:

```

BOOL WriteProcessMemory(
    [in] HANDLE hProcess,
    [in] LPVOID lpBaseAddress,
    [in] LPCVOID lpBuffer,
    [in] SIZE_T nSize,
    [out] SIZE_T *lpNumberOfBytesWritten
);

```

Great, let's set a breakpoint on **WriteProcessMemory** and restart the debugger!
 The 3rd hit is the correct one: we are sure about this since we see from the stack that the *lpBaseAddress* (the address where the data will be written) is right next to our *OEP*:

0012F760	0072F6E5	return to _start.01
0012F764	FFFFFFFF	return to FFFFFFFF
0012F768	0044DBB0	_start.0044DBB0
0012F76C	00CE8050	
0012F770	00000200	
0012F774	0012F76C	

At **0xCE8050** there is the buffer that will be written. Right click on it and select "Follow DWORD in disassembler":

00CE8050	8CE1	mov ecx, fs
00CE8052	46	inc esi
00CE8053	0033	add byte ptr ds:[ebx], dh
00CE8055	D266 8B	shl byte ptr ds:[esi-75], cl
00CE8058	14 41	adc al, 41
00CE805A	83E2 08	and edx, 8
00CE805D	8955 EC	mov dword ptr ss:[ebp-14], edx
00CE8060	837D EC 00	cmp dword ptr ss:[ebp-14], 0
00CE8064	74 0B	je CE8071
00CE8066	8B45 08	mov eax, dword ptr ss:[ebp+8]

We are exactly inside the buffer that will be written replacing the memory starting from **0x44DBB0**. Finding our *OEP* inside the buffer at this point is easy: we can add to the current address (**0xCE8050**) the difference between **0x44DCB0** (address of the *OEP*, found previously) and **0x44DBB0** (base address where the buffer will be written to): **0xCE8150**.

Indeed, finally at this address we find our *OEP* inside the buffer:

00CE814E	CC	int3
00CE814F	CC	int3
00CE8150	55	push ebp
00CE8151	8BEC	mov ebp, esp
00CE8153	6A FF	push FFFFFFFF
00CE8155	68 70454600	push _start.464570
00CE815A	68 D4AE4400	push _start.44AED4
00CE815F	64:A1 00000000	mov eax, dword ptr fs:[0]
00CE8165	50	push eax
00CE8166	64:8925 00000000	mov dword ptr fs:[0], esp
00CE816D	83C4 A4	add esp, FFFFFFFA4
00CE8170	53	push ebx
00CE8171	56	push esi
00CE8172	57	push edi

Great! To halt the execution at the *OEP*, we can modify the buffer, replacing the first two bytes with **EBFE** (infinite loop). In this way the buffer will be written with our patch and once the execution flow reaches the *OEP*, it will get stuck right there in our infinite loop. All that remains at that point is to set a breakpoint to block the execution and restore the original bytes. Let's proceed by patching the following bytes:

00CE814F	CC	int3
00CE8150	EB FE	jmp CE8150
00CE8152	EC	in al,dx
00CE8153	6A FF	push FFFFFFFF
00CE8155	68 70454600	push _start.464570
00CE815A	68 D4AE4400	push _start.44AED4
00CE815F	64:A1 00000000	mov eax,dword ptr fs:[0]
00CE8165	50	push eax
00CE8166	64:8925 00000000	mov dword ptr fs:[0],esp

Just click on RUN and let the SecuROM loader complete its job. We can remove the breakpoint on **WriteProcessMemory** before we run our target executable again, as we don't need it any longer.

Wait a few seconds and then click on PAUSE. We will end up in a loop exactly at the *OEP*:

0044DCAD	CC	int3
0044DCAE	CC	int3
0044DCAF	CC	int3
0044DCB0	EB FE	jmp _start.44DCB0
0044DCB2	EC	in al,dx
0044DCB3	6A FF	push FFFFFFFF
0044DCB5	68 70454600	push _start.464570
0044DCBA	68 D4AE4400	push _start.44AED4
0044DCBF	64:A1 00000000	mov eax,dword ptr fs:[0]
0044DCC5	50	push eax

Let's set a breakpoint at 0x44DCB0 and restore the original opcodes (558B):

0044DCAD	CC	int3
0044DCAE	CC	int3
0044DCAF	CC	int3
0044DCB0	55	push ebp
0044DCB1	8BEC	mov ebp,esp
0044DCB3	6A FF	push FFFFFFFF
0044DCB5	68 70454600	push _start.464570
0044DCBA	68 D4AE4400	push _start.44AED4
0044DCBF	64:A1 00000000	mov eax,dword ptr fs:[0]
0044DCC5	50	push eax

Perfect, the execution flow is currently stuck at the *OEP* ... just like we wanted!

If we now try to dump the binary with Scylla, we will obtain an executable that crashes at the first call (theoretically a **GetVersion**). We need to figure out what SecuROM did to the APIs used by the game (remember: we already found a suspicious call when we used a breakpoint on **GetCommandLineA** earlier).

If you are using a VM, I suggest creating a snapshot of the current state, so after understanding what happens to the APIs, we can quickly restore everything back to the *OEP* without having to start over (a special thanks goes to Antelox for recommending me this technique).

Let's process by stepping some instructions until we get to the first call:

0044DCD1	56	push esi
0044DCD2	57	push edi
0044DCD3	8965 E8	mov dword ptr ss:[ebp-18],esp
0044DCD6	FF15 00137400	call dword ptr ds:[741300]
0044DCDC	A3 5CCE4700	mov dword ptr ds:[47CE5C],eax
0044DCE1	A1 5CCE4700	mov eax,dword ptr ds:[47CE5C]

It would have been reasonable to expect a call to **GetVersion**, but instead we are in front of a call that takes us to a function located in the **.cms_t** segment.

Let's click on step into to enter in this function to study what's happening here.

We realize that we are dealing with a particular function, there are also some calls to **timeGetTime** probably with the aim of detecting if we are spending time stepping between instructions with the debugger.

If we scroll down, we notice that before the classic *RET*, there is a very suspicious instruction, a **jmp eax**:

```

00730270  5F          pop     edi
00730271  5E          pop     esi
00730272  5B          pop     ebx
00730273  8BE5      mov     esp, ebp
00730275  5D          pop     ebp
00730276  FFE0      jmp     eax
00730278  5F          pop     edi
00730279  5E          pop     esi
0073027A  5B          pop     ebx
0073027B  8BE5      mov     esp, ebp
0073027D  5D          pop     ebp

```

This already gives us some clues about what is going on here (especially if you have read my technical paper about Laserlock!) 😊

We put a breakpoint on the jump and press RUN.

Once we hit the Breakpoint, we check the **EAX** register:

```

Hide FPU
EAX  7C81126A  <kernel32.GetVersion>
EBX  7FFD7000  &L"=::=:\"
ECX  00144D30
EDX  7C98B140  ntdll.7C98B140
EBP  0012F7AC
ESP  0012F730

```

Here it is, the call we expected! It will be reached thanks to that jump!

We can run a second test by looking for another call at **0x00741300** immediately after the *OEP*. If you remember, while we were looking for the *OEP* we had set a breakpoint on **GetCommandLineA** and the call where execution was blocked originated from **0x0044DD50**. At that address we have a call dword ptr ds:[741300]:

```

0044DD44  C745 FC 00000000  mov     dword ptr ss:[ebp-4], 0
0044DD4B  E8 207B0000      call   _start.455870
0044DD50  FF15 00137400    call   dword ptr ds:[741300]
0044DD56  A3 18917200      mov     dword ptr ds:[729118], eax
0044DD5B  E8 F0780000      call   _start.455650

```

Let's continue the execution and follow this call. We will obviously come back to the function we have just analysed and in the end, we will find ourselves stuck again on the **jmp eax** (if you have kept the breakpoint active). Now let's check the **EAX** register:

```

Hide FPU
EAX  7C812FAD  <kernel32.GetCommandLineA>
EBX  7FFD7000  &L"=::=:\"
ECX  7FFD7000  &L"=::=:\"

```

Here is the reference to the **GetCommandLineA** API.

Now we no longer have any doubts: SecuROM has replaced the APIs used by the game with its own function (located in the `.cms_t` segment) which, based on the source address of the call, calculates the requested API and reaches it via a jump.

For those who want to understand in detail how the correct APIs are retrieved, you can study the disassembled part enclosed within the *CriticalSection* of that function, paying attention to the various precautions used to make debugging more complicated (such as `timeGetTime`).

At this point it would be reasonable to think of writing a few assembly lines to loop through the entire `.text` segment looking for the SecuROM calls, and after having called them to retrieve the right APIs, patch them with the correct addresses just obtained.

However, there is a big problem: the address stored in `EAX` at the moment of the jump does not pass from the relative thunk in the *IAT*, but instead it is an absolute jump to the requested function! We cannot replace it overwriting the one in the SecuROM call, because it will change (due to being dynamic).

We need the address of the corresponding thunk in the *IAT*.

So, this is the idea:

- 1) Let's loop through the `.text` segment looking for the SecuROM calls
- 2) As soon as we find a SecuROM call we jump into it
- 3) We hook the `jmp eax` instruction, with a jump to get back to our assembly code (getting the direct address of the API we are resolving)
- 4) Using the direct address of the API (stored in `EAX`), we loop through the *IAT* looking for the corresponding thunk
- 5) Once found it, we patch the function in the `.text` segment to call the address of the thunk, defeating SecuROM.

Finding the starting address of the *IAT* is an extremely easy operation: we can go to the relevant section in the **Memory Map** tab, select the `.idata` segment and click on "Follow in dump". Now we will find the first occurrences of addresses to the various APIs used in our target:

```

0072A4B0 A3 CC 6F 73 1B CB 6F 73 00 00 00 00 00 00 00 00
0072A4C0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0072A4D0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0072A4E0 00 00 00 00 26 B1 21 72 00 00 00 00 00 00 00 00
0072A4F0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0072A500 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0072A510 00 00 00 00 3B 47 E8 73 00 00 00 00 00 00 00 00
0072A520 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0072A530 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0072A540 00 00 00 00 C1 61 E4 77 69 5A E4 77 00 00 00 00
0072A550 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0072A560 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0072A570 00 00 00 00 00 00 00 00 BD FD 80 7C FA CA 81 7C
0072A580 EE 94 80 7C BF FC 80 7C 28 1A 80 7C 6B 23 80 7C
0072A590 E2 10 83 7C 12 FF 80 7C AD 2F 81 7C F5 60 83 7C

```

The *IAT* starts at `0x72A4B0` (`0x736FCCA3` is the address of `DirectDrawCreate` in this case). As we can see this *IAT* is a bit peculiar: it is full of free spaces.

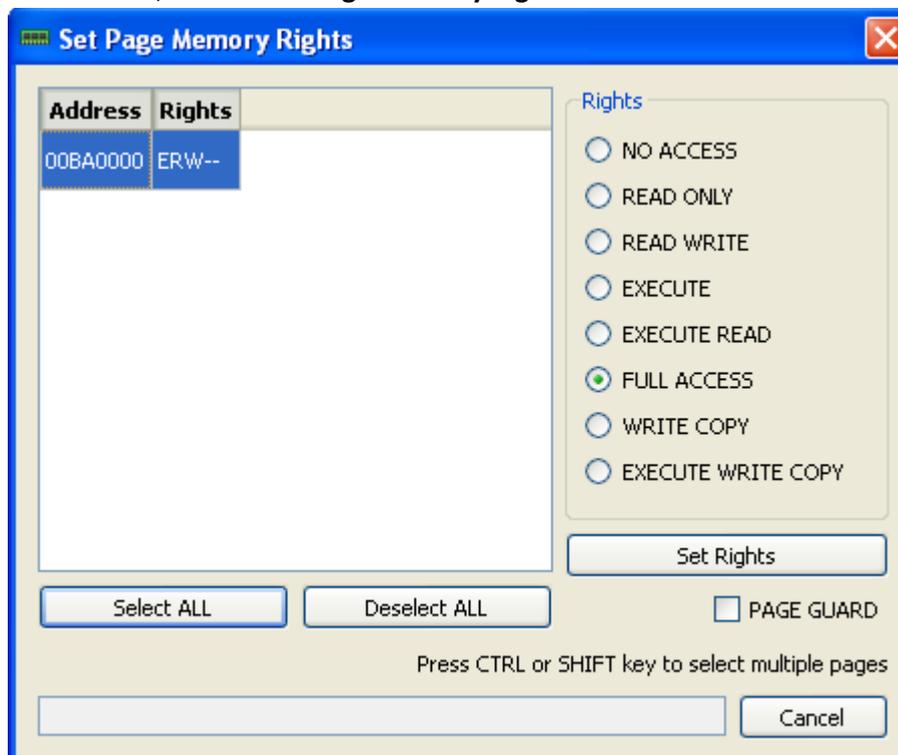
Alternatively, we could have used Scylla to look up the address of the *IAT* for us.

Scrolling down a bit in the dump we can see that the *IAT* ends up at the address **0x72A87F**:

```
0072A7E0 BA 7D B1 76 C1 79 B1 76 2B 84 B1 76 C2 80 B1 76
0072A7F0 92 82 B1 76 AC 7F B1 76 45 AA B1 76 A5 AD B0 76
0072A800 90 B0 B1 76 D4 02 B1 76 F8 94 B0 76 E1 95 B0 76
0072A810 B2 06 B1 76 E1 07 B1 76 E2 43 B1 76 E1 E6 B0 76
0072A820 FC FB B0 76 66 F6 B0 76 B6 5F B0 76 01 52 B0 76
0072A830 F3 05 B1 76 16 01 B1 76 00 00 00 00 00 00 00 00
0072A840 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0072A850 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0072A860 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0072A870 00 00 00 00 00 00 00 00 53 2A 4D 77 7E 05 4D 77
0072A880 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0072A890 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0072A8A0 00 00 00 00 00 00 00 00 00 00 00 00 00 E4 01 4D 75
0072A8B0 6C 74 69 42 79 74 65 54 6F 57 69 64 65 43 68 61
0072A8C0 72 00 64 02 53 65 74 45 72 72 6F 72 4D 6F 64 65
```

The last thing we need is to find a code cave where to put our own assembly code.

From the **MemoryMap** tab we choose a memory area that is marked as **PRV**, that is free and that is large enough for our purpose. I chose the one starting **0xBA0000**. Before continuing, remember to right-click on this section, choose **Set Page Memory Rights** and select **Full Access**:



We confirm by clicking **Set Rights**.

Since we're going to patch the calls in the text segment (replacing the usual SecuROM proxy calls with the real API addresses), let's make sure it has the writeable flag correctly set.

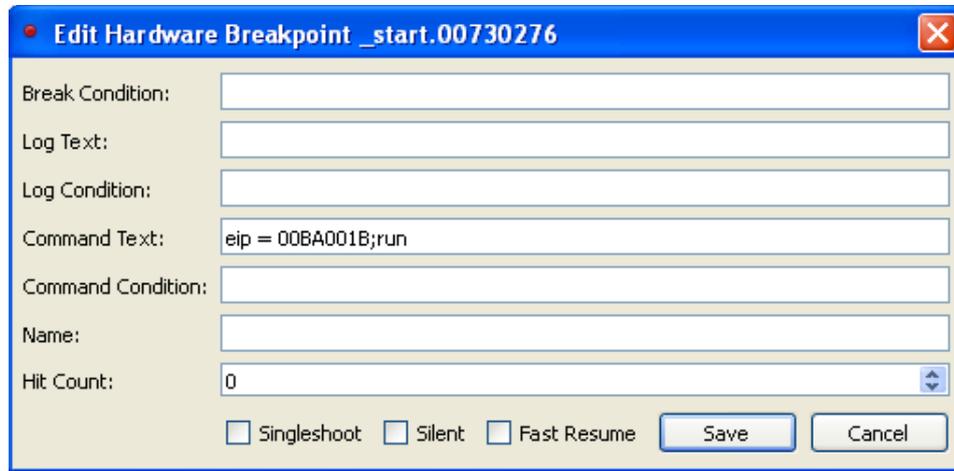
We are ready to write our code now, so let's go to **0xBA0000** and insert the following code:

00BA0000	B9 00104000	mov ecx,_start.401000	
00BA0005	8139 FF150013	cmp dword ptr ds:[ecx],130015FF	
00BA0008	75 2B	jne BA0038	
00BA000D	8079 04 74	cmp byte ptr ds:[ecx+4],74	
00BA0011	75 25	jne BA0038	
00BA0013	890D 9000BA00	mov dword ptr ds:[BA0090],ecx	store ecx
00BA0019	FFE1	jmp ecx	jump to the securom call
00BA001B	8B0D 9000BA00	mov ecx,dword ptr ds:[BA0090]	restore ecx (return from hook)
00BA0021	BB 80A47200	mov ebx,<_start.&DirectDrawCreate>	mov ebx, 0x72A4B0 (IAT START)
00BA0026	3903	cmp dword ptr ds:[ebx],eax	
00BA0028	74 0B	je BA0035	
00BA002A	43	inc ebx	
00BA002B	81FB 7FA87200	cmp ebx,_start.72A87F	
00BA0031	75 F3	jne BA0026	
00BA0033	CD 03	int 3	ERROR!! THINK NOT FOUND!!
00BA0035	8959 02	mov dword ptr ds:[ecx+2],ebx	
00BA0038	41	inc ecx	
00BA0039	81F9 F92F4600	cmp ecx,_start.462FF9	
00BA003F	75 C4	jne BA0005	
00BA0041	CD 03	int 3	COMPLETED!
00BA0043	0000	add byte ptr ds:[eax].a1	

This is what the above code actually does:

- First it puts the starting address of the text section into the **ECX** register. The bytes are then compared against **0xFF15001374**, i.e., with the call to the SecuROM function (performed in two separate **cmp** instructions to check all the related bytes): If there is no match the **jne branch** will be followed and the address in **ECX** will be incremented by 1 to check for the next byte.
- If, on the other hand, we are in the front of a SecuROM call, we save the current address in the **ECX** register (which tells us which byte of the **.text** section we have arrived at) and jump into it.
- Once we reach the now famous **jmp eax** located in the SecuROM function we will set a hook (in a few minutes) to automatically jump back to **0xBA001B**. At this point we have the direct address of the requested API inside the **EAX** register.
- Now we restore the **ECX** register and load the start address of the **IAT** into the **EBX** register.
- Let's loop through the **IAT** until we find the thunk that points to the address contained in **EAX** (**EBX** will obviously be incremented each time: If none of the thunks contain the API we're trying to fix, then we're in big trouble (**INT 3** at **0xBA0033**), but obviously this should NOT happen.
- If the thunk has been found, we will jump to address **0xBA0035**, where we will replace the bytes related to the SecuROM call with those of the correct thunk. At this point we can go back to continue scrolling the **.text** segment looking for the remaining calls.
- When the **.text** segment ends (**ECX** will be **0x462FF9**, i.e., the last address of the segment -6, which is the size of the call, in bytes), we are done and can proceed with the dump!

Before running the code, remember to right-click on the line **BA0000** and choose “Set New Origin Here”. Now let's set our hook: move to **0x730276** (the address where the **jmp eax** is located), right click on it and choose “Breakpoint” -> “Set Hardware on Execution”. Let's move to the **Breakpoints** tab, right click on the hardware breakpoint we just created and choose Edit. Let's configure it in this way:



This way once the BP is triggered, we will automatically return to our code (at address 0xBA001B).

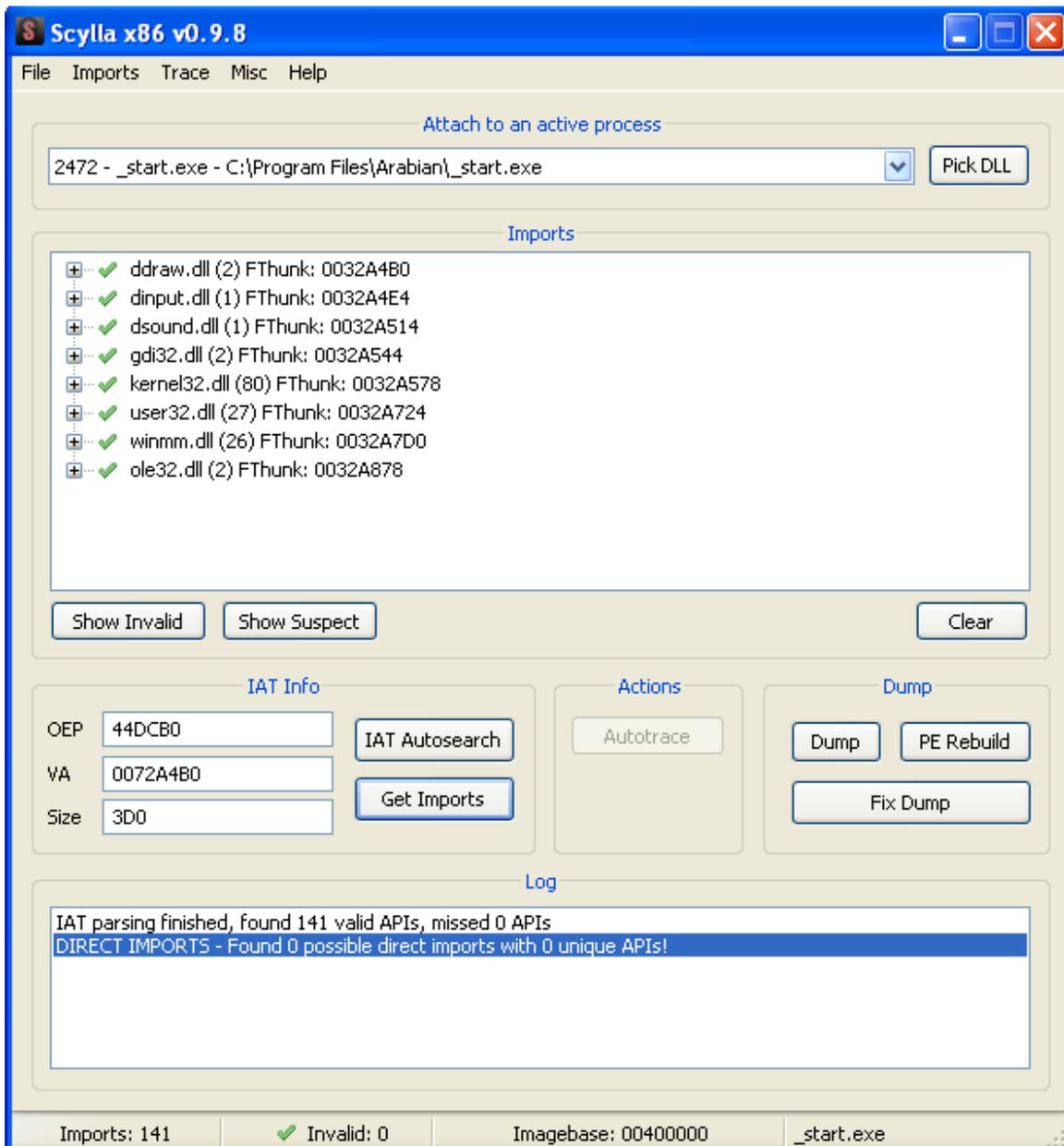
PAY ATTENTION: only use a hardware breakpoint in this case, otherwise the program will crash (the presence of software breakpoints will be detected).

We are ready to run our code, move to **0xBA0000** and click **RUN**.

Once the execution is complete, we will be stuck at **0xBA0041**:

00BA0000	B9 00104000	mov ecx,_start.401000	
00BA0005	8139 FF150013	cmp dword ptr ds:[ecx],130015FF	
00BA0008	75 2B	jne BA0038	
00BA000D	8079 04 74	cmp byte ptr ds:[ecx+4],74	
00BA0011	75 25	jne BA0038	
00BA0013	890D 9000BA00	mov dword ptr ds:[BA0090],ecx	store ecx
00BA0019	FFE1	jmp ecx	jump to the secur
00BA001B	8B0D 9000BA00	mov ecx,dword ptr ds:[BA0090]	restore ecx (ret
00BA0021	BB B0A47200	mov ebx,<_start.&DirectDrawCreate>	mov ebx, 0x72A4E
00BA0026	3903	cmp dword ptr ds:[ebx],eax	
00BA0028	74 0B	je BA0035	
00BA002A	43	inc ebx	
00BA002B	81FB 7FA87200	cmp ebx,_start.72A87F	
00BA0031	75 F3	jne BA0026	
00BA0033	CD 03	int 3	ERROR!! THINK NC
00BA0035	8959 02	mov dword ptr ds:[ecx+2],ebx	
00BA0038	41	inc ecx	
00BA0039	81F9 F92F4600	cmp ecx,_start.462FF9	
00BA003F	75 C4	jne BA0005	
00BA0041	CD 03	int 3	COMPLETED!
00BA0043	0000	add byte ptr ds:[eax],a1	
00BA0045	0000	add byte ptr ds:[eax],a1	
00BA0047	0000	add byte ptr ds:[eax],a1	
00BA0049	0000	add byte ptr ds:[eax],a1	

We are almost there! Let's launch Scylla, choose the right process (**_start.exe**), set the **OEP**, the address of the **IAT** and its size:



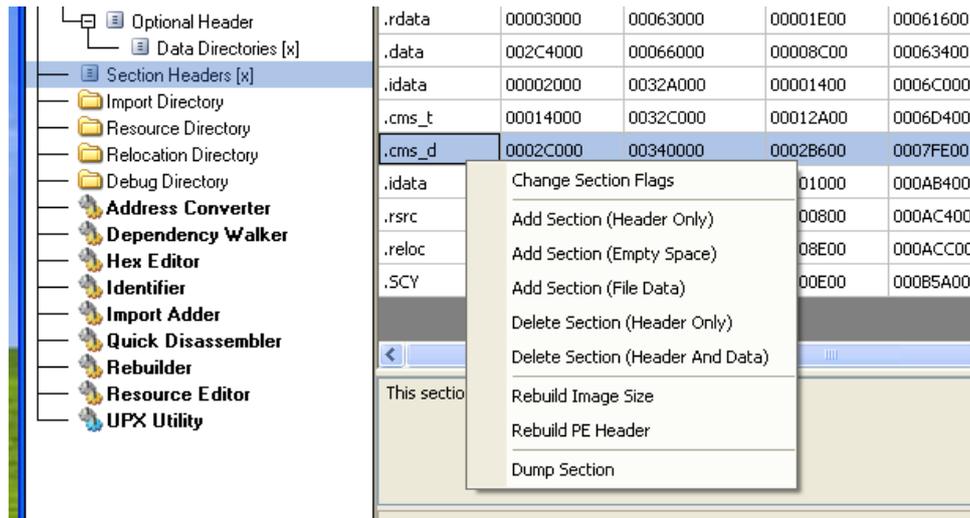
Let's click on *Get Imports*, then on *Dump* and finally on *Fix Dump*.

Finally, we will have an executable free from SecuROM 😊

Let's finish the job:

Our executable works perfectly and SecuROM is just past the memory addresses used by our new executable. However, if we want to be true perfectionists, we can make it slightly smaller by removing those two sections used by the SecuROM loader that now only take up useless space.

Let's load our binary into CFF Explorer and move to the *Section Headers* tab. Select the *.cms_t* and *.cms_d* sections, right click and choose *Delete Section (Header And Data)*:



Let's save the new executable (it will be much smaller) and we're **DONE!** 😁

Credits:

I would like to thank again the legendary Antelox for suggesting me the snapshot technique on VM. Undoubtedly a great way to save a lot of time.

A big thank you goes to m00k00 for reviewing and fixing the english translation of this technical paper!!

Conclusion:

As we have seen SecuROM *new* 4.48.00.004 conceptually shares something with Laserlock. Ultimately it is a very didactic DRM from which we have certainly learned something.

Thank you for reading this document 😊

Luca